A Case for Standard Non-Blocking Collective Operations

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Collective Operations

S. Gorlatch: "Send Receive considered harmful"

"Parallel programming based on message passing can be improved by expressing communication in a structured manner without using send-receive statements."

Collective Operations

- provide common communication patterns
- improve performance (portability)
- programability and code maintenance (readability)
- avoid implementation errors

Overlapping Communication and Computation

J. Sancho: "Quantifying the Potential Benefit of Overlapping Communication and Computation in Large-Scale Scientific Applications"

"This result indicates that for a potentially large class of real-world applications, future high performance computing systems could use lower performance and more cost-effective networks without negatively impacting application performance if the overlap of communication with computation were fully exploited."

Overlap of Communication and Computation

- hides network latency
- enables the use of cheaper networking technology
- increases the tolerance of process skew
- collective communication overlap?

Non-blocking Collective Operations

R. Brightwell: "Implications of Application Usage Characteristics for Collective Communication Offload"

"Even so, current code shows the opportunity to overlap from thousands of instructions (several microseconds) to hundreds of thousands of instructions (approaching a millisecond) with many of the collective calls."

Non-blocking collectives

- combine many advantages
- are not available in the current MPI-2.0 standard
- begin to gain more and more attention
- missed by scientists who replace collective patterns with non-blocking point-to-point operations (cf. Danalis et al. "Transformations to Parallel Codes for Communication-Computation Overlap")

Our Standard Proposal

- we define an API similar to MPI
- relaxes the MPI semantics (allows more than one collective per communicator)
- no tags (ordering is crucial for matching)
- no defined asynchronous progress as we are used to ;-)

Programming Interface Example

```
MPI_Request request;
MPI_Ibarrier(comm, &request);
...
/* computation, other MPI communications */
...
MPI_Wait(request, status);
```

LibNBC - Non-blocking Collectives for MPI-2

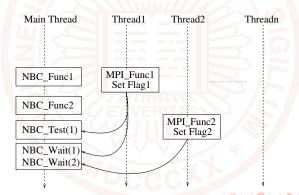
- open-source tuned implementation of asynchronous collectives
- two implementation options evaluated:
 - threads (recommended by MPI Forum)
 - point-to-point message based implementation

Programming Interface Example

```
NBC_Handle request;
NBC_Ibarrier(comm, &request);
...
/* computation, other MPI communications */
...
ret = NBC_Wait(request);
```

Implementation Option 1 - Threads

- needs MPI THREAD MULTIPLE
- task queue model
- occupies one CPU per process for communication
- MPI COMM DUP problems (Gropp et al. "Issues in Developing a Thread-Safe MPI Implementation")



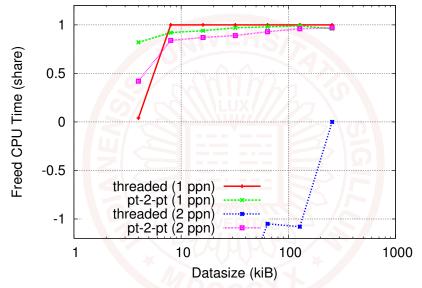
Implementation Option 2 - Point-to-point

- needs only MPI-1 functions (MPI-2 for Fortran bindings)
- schedule-based execution
- performance depends on overlap of ptp messages
- implementation of the fastest known algorithms (static selection)
- available since 2006, first users adapt their codes
- no matching with MPI collective functions

| recv from 0 end send to 3 | end | send to 5 | |
|---------------------------|-----|-----------|--|
|---------------------------|-----|-----------|--|

Figure: Binomial broadcast schedule on rank 1 of 6.

Microbenchmarks - Freed CPU Time



64/128 processes on 64 Coyote System nodes (Dual Opteron, IB)

Non-blocking 3D-FFT

Derivation from "normal" implementation

- distribution identical to "normal" 3D-FFT
- first FFT in z direction and index-swap identical

Design Goals to Minimize Communication Overhead

- start communication as early as possible
- achieve maximum overlap time

Solution

- start MPI_Ialltoall as soon as first xz-plane is ready
- calculate next xz-plane
- start next communication accordingly ...
- collect multiple xz-planes (tile factor)



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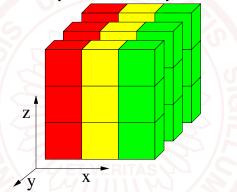
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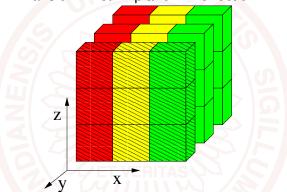


Data already transformed in y direction



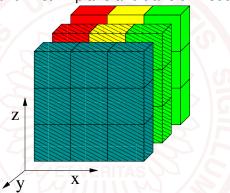
1 block = 1 double value (3x3x3 grid)

Transform first xz plane in z direction



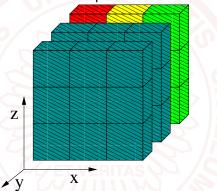
pattern means that data was transformed in y and z direction

start MPI_Ialltoall of first xz plane and transform second plane



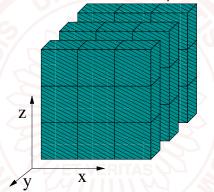
cyan color means that data is communicated in the background

start MPI_Ialltoall of second xz plane and transform third plane



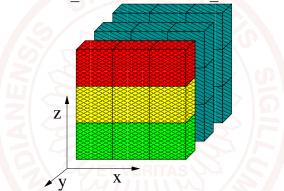
data of two planes is not accessible due to communication

start communication of the third plane and ...

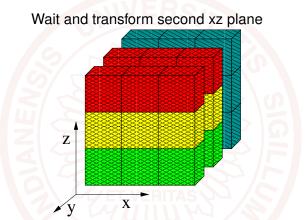


we need the first xz plane to go on ...

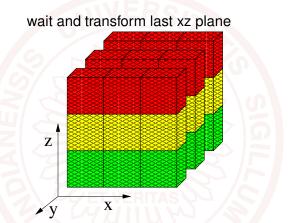
... so MPI_Wait for the first MPI_lalltoall!



and transform first plane (new pattern means xyz transformed)

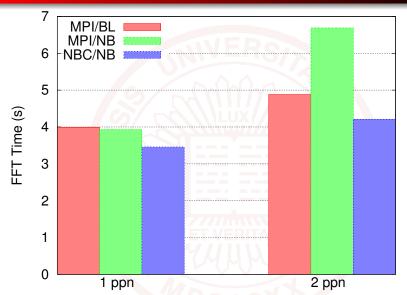


first plane's data could be accessed for next operation



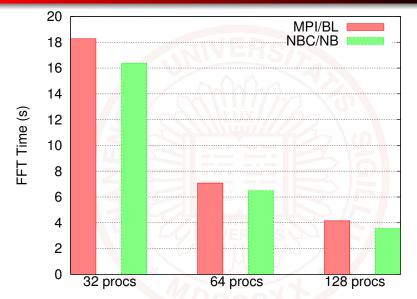
done! → 1 complete 1D-FFT overlaps a communication

1024³ 3d-FFT over InfiniBand



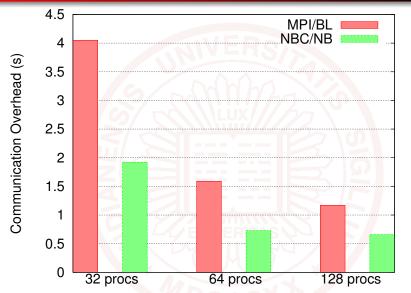
P=128, "Coyote"@LANL - 128/64 dual socket 2.6GHz Opteron node

1024³ 3d-FFT on the XT4



• "Jaguar"@ORNL - Cray XT4, dual socket dual core 2.6GHz Opteron

1024³ 3d-FFT on the XT4 (Communication Overhead)



• "Jaguar"@ORNL - Cray XT4, dual socket dual core 2.6GHz Opteron

It works, why do we need it standardized?



- enable hardware-optimized implementations
- solve software technological dilemma of double implementation
- enable matching of blocking and non-blocking collectives
- use MPI test and wait functions
- overcome problems with first call and Comm_dup()
- → widen user-base



Conclusions and Future Work

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- a threaded implementation "wastes" cores
 - may be acceptable in the future
 - not on today's dual/tri/quad core systems
 - ⇒ LibNBC is currently implemented on top of pt2pt
- standardization would solve many problems!



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Future Work:

- tuning LibNBC and adding it to Open MPI
- port applications to use NBC
- ⇒ We would like to collaborate with scientists!

LibNBC Download/Further Information

http://www.unixer.de/research/nbcoll/



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